# PERFORMANCE STUDY OF AD HOC WIRELESS LAN FOR MULTIMEDIA STREAMING APPLICATIONS

By

## KON PEI SIN

## FINAL REPORT

Submitted to the Electrical & Electronics Engineering Programme in Partial Fulfillment of the Requirements for the Degree Bachelor of Engineering (Hons) (Electrical & Electronics Engineering)

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## **CERTIFICATION OF APPROVAL**

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A project dissertation submitted to the Electrical & Electronics Engineering Programme Universiti Teknologi PETRONAS in partial fulfillment of the requirement for the Bachelor of Engineering (Hons) (Electrical & Electronics Engineering)

Approved by:

Ms. Nasreen Badruddin Project Supervisor

## UNIVERSITI TEKNOLOGI PETRONAS TRONOH, PERAK

June 2006

## **CERTIFICATION OF ORIGINALITY**

This is to certify that I am responsible for the work submitted in this project, that the original work is my own except as specified in the references and acknowledgements, and that the original work contained herein have not been undertaken or done by unspecified sources or persons.

Kon Pei Sin

## ABSTRACT

Wireless Local Area Network (WLAN) provides high-speed and cable-free access for computer-to-computer information transfer within the coverage range. The optional feature of WLAN is ad hoc, which allows the network to be formed and deformed without the need of system administration or access point. Multimedia streaming through WLAN allows users to download and play a small chunk of the multimedia file simultaneously. An ad hoc base wireless network is used in this project rather than infrastructure base due to the simpler architecture, cheaper set up cost of ad hoc network and no system administrator required. Unlike file sharing, streaming will not download the whole file. Streaming can be used for multimedia distribution without users having the illegal copies of the file. In this project, a multimedia file is to be streamed, but not downloaded, through the WLAN without passing through any access point. The host computer broadcast the multimedia file using "media server configurator" while the client computer streamed from the host computer using "streaming media player". Multihop streaming allows a node that is totally out of the coverage range to stream from the host computer, using the bridge as a passage. The streaming of the multimedia file is seamless and smooth within the ad hoc network. A node that is out of the coverage range can stream from the host computer by using the multi-hop concept although the distance between the host and the client does affect the performance. In multi-hop wireless network, the streaming process is at its best when the bridge is about 3 meters away from the receiving node.

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## CHAPTER 1 INTRODUCTION

#### 1.1 Background Study

Wireless Local Area Networks (WLAN) has played a key role in revolutionizing the use of technology in society. The term wireless refers to the transmission of information and data without wires. WLAN provide high-speed, cable-free access for computer-to-computer information transfer within the coverage range. Most installed wireless LAN today is infrastructure-based, which requires the use of access point. The access point provides an interface to a distribution system, which enables wireless users to utilize corporate servers and Internet applications. [1]

The optional feature of wireless LAN standard, 802.11b/g specifies the "ad hoc" mode, which allows the network interface card (NIC) to operate in an independent basic service set (IBSS) network configuration. No access points are needed in the IBSS. [2] An ad hoc wireless network is a collection of two or more devices equipped with wireless communications and networking capability. Such devices can communicate with another node that is immediately within their coverage range. An ad hoc wireless network is self-organizing and adaptive. This means that a formed network can be deformed without the need of system administration. The term "ad hoc" tends to imply "can take different forms" and "can be mobile, standalone, or networked". [3]

Peer-to-Peer (P2P) Networking feature provided by most of the operating system (OS) such as Windows, Linux, and many other OS has a very simple configuration to set up an Ad Hoc network. Two or more PCs can form a peer-to-peer network. In Windows, the P2P grouping is called a workgroup.

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Streaming is a process where users download and play a small chunk of the multimedia file smoothly at the same time but without having a copy of the original file. Streaming is different from file sharing where it does not download the whole file. It is normally used for multimedia distribution without the users having illegal copies of the file. It allows the owner of the file to share the contents but the users would not be able to own the file. The live streaming is very useful in broadcasting the football match, videoconference, net meeting, and deliver lecture. Capture devices record and connect to the host computer for broadcast the captured files while the client computer streamed from the host computer as shown in Figure 1 below.

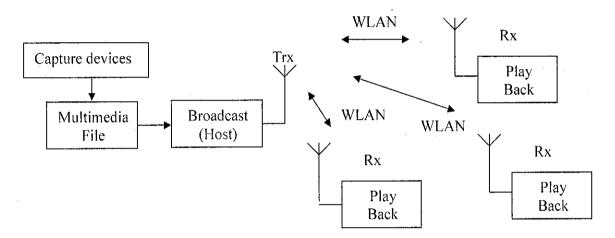


Figure 1: Multimedia Streaming / Live Streaming through WLAN

#### 1.2 Problem Statement

The file sharing and downloading files in Local Area Network (LAN) are very common. In file sharing, the multimedia file shared on the network needs to be completely downloaded in order to play the file. The users will never know the content of the multimedia files until they completely download the file and play it. Thus, the users might not have enough disk space to store such a big multimedia file.

Streaming only allows users to play the multimedia file while downloading small chunks of the file from the source. It will continue download the small chunks while it is playing the multimedia file smoothly. Streaming does not allow the users to download the whole file. It can be used for multimedia distribution without users having illegal copies of the original multimedia file. It allows the owner of the file to share the contents but the user would not be able to own the file.

The basic architecture of networking must be studied in order to proceed with the project. The node discovery in a WLAN, the IP address assign, and the routing of the ad hoc network are the problems that need to be defined. The most challenging part of this project is to allow the node, which totally out of the coverage range to stream the multimedia files by using the multi-hop concept.

#### 1.3 Objectives and Scope of Study

The project goals are as follows:

- Create an ad-hoc wireless network
- Any node within the coverage range should be able to detect the network created
- Users are only allowed to stream the multimedia files but not download the files from the Ad-Hoc network.
- The playing of the multimedia files should be seamless and smooth.
- Create a multi-hop wireless network
- Streaming process in the multi-hop wireless network.

The main purpose of this project is to stream the multimedia file between the host computer and the others node computers with wireless LAN enabled without passing through access point. The streaming process protects the originality of the multimedia files without the users having a copy of the multimedia file. The secondary objective is to learn how to construct an Ad Hoc wireless network in a room to allow multimedia file streaming in the network. The Ad Hoc wireless network created is able to detect any node within the coverage range of the network. The project is being done in Ad-Hoc base rather than the infrastructure base (access point) because the ad hoc network is more suitable for small area data communication and no system administrator is needed. The complexity and cost of setting the

network is lower compared to infrastructure base. The multi-hop streaming is one of the objectives, which allow a node that is totally out of the coverage range to stream the multimedia file from the host.

The project concentrates on the creation of wireless network, node discovery in wireless local area network (WLAN), streaming between the two or more wireless enabled computers, and multi-hop streaming. The Ad-Hoc wireless networking, WLAN, multi-hop networking, and the performance of the multi-hop networking are the issues included in the scope of this study.

## CHAPTER 2 LITERATURE REVIEW / THEORY

### 2.1 Wireless Local Area Network (WLAN)

Wireless Local Area Network (WLAN) provides high-speed, cable-free access for computer-to-computer information transfer within the coverage range. They possess all the functionality of local area network (LAN), but without the physical constraints of the wire itself. The wireless nature allows easy implementation, portability, and lower installation costs. Wireless systems can be installed in different environments, such as offices, manufacturing floors, hospitals, or universities.

Wireless networks have their own standards as in wired network. In 1990, the IEEE (Institute of Electrical and Electronic Engineers) formed a group to develop a standard for wireless equipment. The applicable standard for most of the WLAN world is the IEEE 802.11 WLAN standard. The original 802.11b standard limited WLAN speed to 2 Mbps, which is not efficient enough for most networking needs. In September of 1999, The B standard offered two higher speeds, 5.5 Mbps and 11 Mbps. At the same time B was released, A was also added to the 802.11 standard. 802.11a has a maximum speed of 54 Mbps and supports speeds of 48, 36, 24, 18, 12, 9, and 6 Mbps at a frequency of 5 GHz [4]. The reason that 802.11a is faster than 802.11b is because of its higher frequency, more transmission channels, multiplexed transmissions, and an efficient error-checking scheme. In late 2001, a draft was proposed called 802.11g. 802.11g runs on the same RF band as B (2.4 GHz) but uses the transmission techniques of A [4].

There are some challenges faced on WLAN such as radio signal interference. It is possible to have third party interference that using wireless bridging or other wireless devices operating in the same frequency band. Many other devices, such as cordless phone (portable phone), microwave ovens, wireless speakers, blue-tooth devices, and security devices, also use these frequencies. Power consumption is always an issue with laptops and PDAs, because the power and battery have limited lives. In a closed network, a single vendor solution has advantage on interoperability and management to avoid the possibility of one vendor blaming the other for equipment failures. Security in the IEE 802.11 specification, which applies to 802.11b/a/g, has come under such analysis. IEEE enhanced Wired Equivalent Privacy (WEP) to provide robust authentication options with 802.1x to make 802.11 based WLAN secure. At the same time, IEEE is looking to stronger encryption mechanisms [5]. WLAN using the TCP/IP protocols helps protect the network against any loss or corruption of data over the air. Because of differences in component configuration, placement, and physical environment, every infrastructure application is a unique installation. There are many unknowns concerning the safe limits of human exposure to RF radiation [5].

### 2.2 Ad-Hoc Wireless Network

An ad hoc network is a collection of wireless mobile nodes dynamically forming a temporary network without the use of any existing network infrastructure or centralized administration. A mobile ad hoc network (MANET) is a system of wireless mobile nodes, which can dynamically self-organize into temporary network topologies. MANET allows people and devices to inter-network in areas without any existing of infrastructure (access point). The formed network can be deformed on the fly without the need for any system administration.

Ad hoc nodes should be able to detect the presence of other such devices and to perform the necessary handshaking to allow communications and sharing of the information and services [6]. Ad hoc devices not only can detect the presence of connectivity with neighbouring devices/nodes, but also identify what type the devices are and their corresponding attributes. Since an ad hoc wireless network does not rely on any fixed network entities, the network itself is essentially infrastructure-less. There is no need for any fixed radio base stations, any wires or fixed routers. However, due to the presence of mobility, routing information will have to change to reflect changes in link connectivity [6].

There are some challenges faced in mobile ad hoc networks. To prevent interference, ad hoc networks must operate over some form of allowed or specified spectrum range. Most microwave ovens operate in the 2.4GHz band, which can therefore interfere with wireless LAN systems. Since multiple mobile ad hoc nodes share the same media, access to the common channel must be made in a distributed fashion, through the presence of a MAC protocol. The MAC protocol must compete for access to the channel while at the same time avoiding possible collisions with neighbouring nodes [6]. Ad hoc mobile networks are different from packet radio networks since nodes can move freely resulting in a dynamically changing topology. In ad hoc mobile networks, mobile devices must perform both the role of an end system and an intermediate system. Hence, forwarding packets on the behalf of others will consume power, and this can be quite significant for nodes in ad hoc wireless network. Ad hoc networks are intranets and they remain as intranets unless there is connectivity to the Internet. Such confined communications have already isolated attackers who are not local in the area.

To join an IP network and communicate with others, a node needs to be configured either manually by an administrator or automatically through a server. However, the manually configured is impractical for large networks as well as in the case of a mobile ad hoc network due to mobility of nodes. There are two ways to distribute IP address auto-configuration mechanisms for mobile ad hoc networks, which is RADA (Random Address Allocation) and LiA (Linear Address Allocation). RADA is base on random IP address selection, while LiA linearly assigns new addresses by utilizing the current maximum IP address value. There is an improved version of LiA, known as LiACR (Linear Allocation with Collision Resolution), which reduces control overhead.

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#### 2.3 Multi-hop Wireless Network

Multi-hop wireless networks combine the characteristics of both wireless and mobile ad hoc networks. In multi-hop wireless networks, the existence of a communication link between the node and the host is not required. Indeed, a (usually small) number of other nodes relay the communication between the node and the host.

Several benefits can be expected from the use of multi-hop wireless network. First, the energy consumption of the mobile devices (nodes) can be reduced. The energy consumption required for radio transmission grows linearly with the distance. Therefore, the battery life of wireless devices can be extended if packets are routed in small hops from the host to the node. Second, as an immediate positive side effect of the reduced transmission energy, the interference is reduced. Third, if the nodes are not too remote from each other, they can communicate independently from the infrastructure. Fourth, the number of fixed antennas can be reduced. Finally, the coverage of the network can be increased using multi-hop wireless network [7].

Although Multi-hop wireless networks are attractive at first sight, multi-hop wireless networks raise a number of problems. For example, in conventional wireless networks with access point, base stations usually are in charge of channel allocation, the synchronization, and power control of mobile devices. To accomplish these tasks, the base stations take advantage of their direct communication link with every mobile device currently visiting their cell. It is quite difficult to extend these operating principles to multi-hop wireless networks.

### 2.4 Peer-to-Peer (P2P) in Windows

Fortunately, Windows, Linux, and many other OS provide for a very simple ad hoc network called peer-to-peer networking. In a conventional computer network, a server can support many PCs. The server and PC are very different, in both software and function. The server functions are to authenticate PS users, interconnect PCs to the server and each other. This architecture is called client/server. The individually networked PCs act as clients, rather than peers of the server. However, there is another method to connect PCs into network [8]. Two or more PCs can form a peerto-peer network. In Windows, this P2P grouping is called a workgroup.

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# CHAPTER 3 METHODOLOGY / PROJECT WORK

## 3.1 Procedure Identification

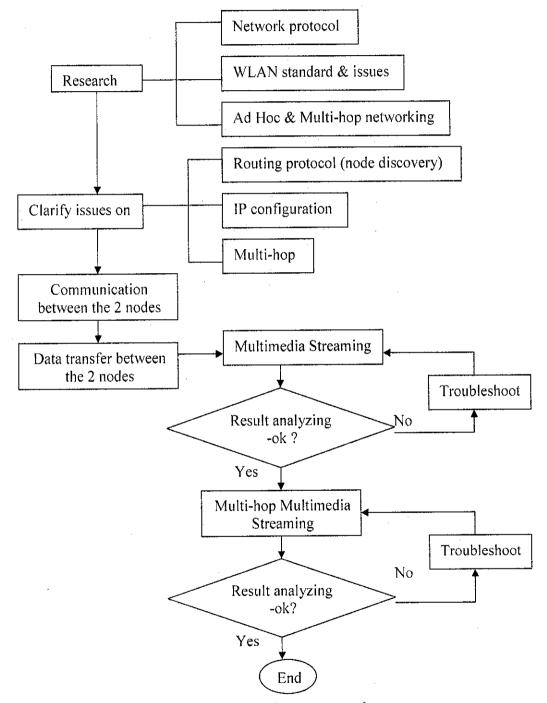


Figure 2: Final Year Project planning procedures

The planning procedure for the final year project is clearly illustrated in Figure 2. In the earlier stage of the project, research on the network protocol, wireless local area network (WLAN), the ad-hoc wireless network, and the multi-hop wireless network has been done. Then, the issues on routing protocol, IP configuration, and multi-hop networking are clarified. The project is continued by conducting the wireless network within two nodes in a laboratory at block 22. After the network set up, the data are transferred from one node to the other to prove the wireless network is functioning. Multimedia file streaming is experimented by using the network created with different distance between the nodes. The project proceeds with the multi-hop multimedia file streaming. Another node, which is totally out of the coverage range, can stream from the host by using the multi-hop concept. Experimental on the distance between nodes and the data transfer rate of the multi-hop streaming is carried out to evaluate the performance of the streaming in multi-hop wireless network.

## 3.2 Tools / Equipments Required

 Two wireless LAN enabled computers Nowadays, all laptops come with the integrated Wi-Fi enable card with the latest technology of Intel Centrino<sup>®</sup>

#### 2. 802.11b/g USB wireless adapter

A USB 2.0 wireless adapter, which is used to attach to the desktop computer so that it becomes WLAN enabled. This USB wireless adapter uses the IEEE 802.11g standard

### 3. Multimedia file

A video clip, MTV, and anime, which normally for entertainment purposes are considered as a multimedia files.

### 4. Media server configurator

The software used at the host computer to broadcast the multimedia files

5. Streaming media player

The software used at the client computer to stream the multimedia file from the host computer.

6. Ethereal

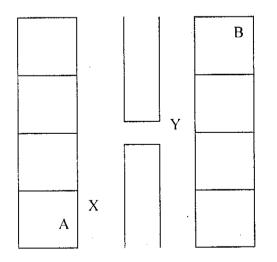
The software used to evaluate the packet / data transfer rate in a wireless network.

7. Measuring tape

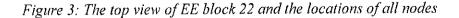
A measuring tool to measure the distance between nodes

#### 3.3 The Experiment Environment

All the experiments and troubleshooting of this project is done in the laboratories in block 22. The set up of the wireless network between nodes A, B, X, and Y are shown in Figure 3 where node A and node B are desktop with the USB wireless adapter while node X and node Y are laptop with wireless enabled (Intel Centrino<sup>®</sup>).

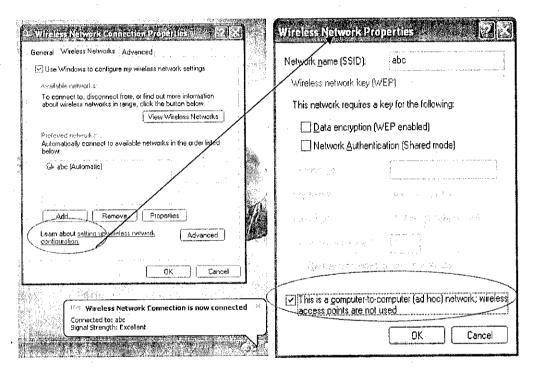


\*The drawing is not according to scale



### 3.4 Setting Up The Wireless Local Area Network (WLAN)

An experiment was conducted to detect the Wireless network created by the 54 Mbps 802.11g wireless USB 2.0 adapters. A wireless local area network (WLAN) is created using the laptop (Intel Centrino<sup>®</sup>) with ad-hoc peer-to-peer connection. The wireless local area network is set up by simply give a network name (SSID) and configuring the network authentication to "open" and data encryption to "disable" so that no network key is required. The settings of the wireless local area network are as shown in Figure 4.



*Figure 4: The setting of a wireless local area network.* 

To ensure that the computer does not connect to any access point within the coverage range, the option "This is a computer-to-computer (ad-hoc) network; wireless access points are not used" at the bottom must be checked as shown in Figure 4. The setting is completed by clicking the "OK" and an icon pop-up at the right bottom corner indicated the wireless network connection is now connected.

After the ad-hoc peer-to-peer wireless network is set up, other nodes, which are within the coverage range, can easily connect to the wireless network created. The

desktop with the wireless USB adapter or any Intel Centrino<sup>®</sup> laptop shows the wireless network icon on the right bottom corner as shown in Figure 5.

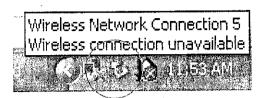


Figure 5: The icon of wireless network connection

The wireless network connection icon is double clicked. The window as in Figure 6 and Figure 7 shows the available network connection within the coverage range for the desktop with USB wireless adapter and laptop with Intel Centrino<sup>®</sup> separately. The specific network name is selected, thus, connection is made to the particular network by clicking on the connect button.

Available netv	vorks:
¦ mywlan ⊘glydd ⊘ FYP	
	n an an an an Arthreachailte an Arthreachailte an an Arthreachailte an Arthreachailt
	2010-201 - 100-00 - 100-00 - 100-00 - 100-000-0

Figure 6: The Wireless network connection available (desktop - USB wireless adapter)

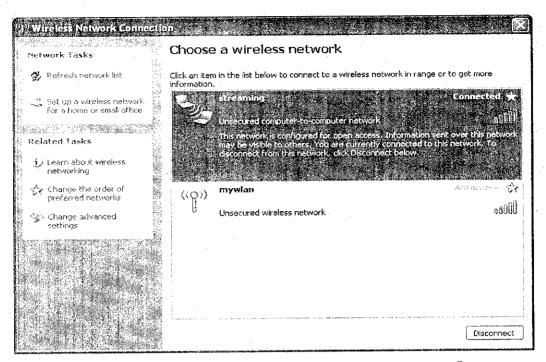


Figure 7: The Wireless network connection available (Intel Centrino<sup>®</sup> laptop)

After the devices are connected to the wireless network, it a pop up message appears on the right bottom corner to indicate the wireless network is connected. It is found that both nodes connected to the wireless network are automatically assigned with individual IP address as shown in Figure 8.

General Support	Wireless Network Connection 3 Status
Address Type: Automatic Private Address	internet Protocol (TCP/IP) Address Tune: Automatic Private Address
IP Address: 159.254.141.176 Subnet Mask: 255.255.00 Default Gateway:	Address Type: Automatic Private Address IP Address: 169.254.62.68 Subnet Mask: 255.255.0.0
Utails           Windows did not detect problems with this connection. If you cannot connect, click	Default Gateway: Details
Flepair.	Regair
Close	<u>[lose</u> ]

Figure 8: The Wireless network connection status of both nodes

#### 3.5 Streaming Process

The general streaming process is shown in Figure 9. The wireless network is set up accordingly as mention in section 3.4. The file sharing and downloading is to ensure the stability of the wireless network created. The host computer as well as the client computer is installed with the software "media server configurator" and "streaming media player" separately, are configured accordingly. The host computer IP address is entered and the streaming process is initiated.

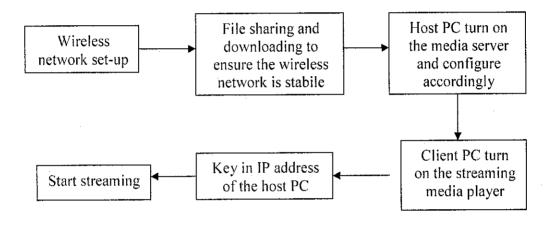


Figure 9: Steaming process procedures

### 3.5.1 Setting Up The Host Computer (Server)

The software "media server configurator" is installed in the host computer. To begin the streaming process, the software - media server configurator is started and a new virtual folder is added as shown in Figure 10. The virtual folder is added by browsing through the folder in the hard disk drive of the host computer where the multimedia file is stored. A description and folder name is required for the multimedia files folder. The other settings are left as the default options, as shown in Figure 11.

Help			Accessibility	Description
Media Server     Media Server     Media Server     MediaR     MediaR     N	Name In (MediaDoct Iew virtual Folder	Path C:\Program Files\U	Anonymous access	Description
	ey. A La gradje veza na			
		nga-ung panjagan karangan na pangan pang		10000 0400000 0400000000000000000000000
	Active resources: 1	Vame		Number of users
		,		

*Figure 10: The media server configurator* 

Folder name	media		ок
Browse	C:\Documents and Sett	ings\Administrator\My	Cancel
Description	try	1. มี และกับสารางการี แบบจะที่มีขับปฏิบัติเราร้องจะมีมีการหรือ เหมือสามารี กับและจะสำนักที่ได้เสียง	n dolandin da da di ama ny yenika an'ny ana ana bash
🕅 Unicast d	elivery 🔽 Multicast deliv	very Group 225.0.0.3:	5003
🖗 Allow to t	eat folder as a Playlist	Order 🌾 Alphabet	
	eat folder as a Playlist	·	
-Authenticati	eat folder as a Playlist	Order 🔅 Alphabet	
-Authenticati	eat folder as a Playlist on method lefault settings nous access)	Order 🏾 🏵 Alphabet	

Figure 11: The virtual folder set-up

When the setting of the virtual folder is done, it shows the list of multimedia files in the particular location of the hard disk drive as in Figure 12. By right clicking on the folder name and choosing "start multicasting", the server starts to broadcast the multimedia files to the entire wireless network.

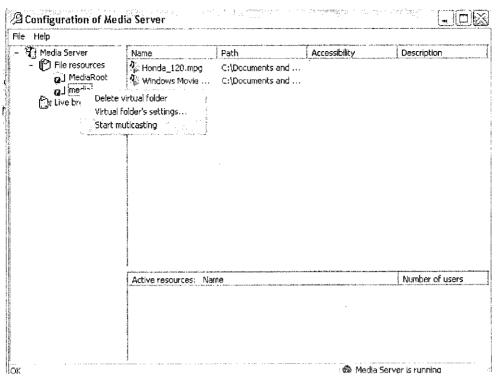


Figure 12: Broadcast the multimedia files from the host computer

## 3.5.2 Setting Up The Client Computers (nodes)

The client computer is installed with the software "streaming media player" in order to stream the multimedia files from the host computer. The streaming media player is started. In the "play" tab on the left top corner, the *Play remote playlist* is clicked and a window as in Figure 13 appears. The IP address of the media server is entered in the column provided. The other important information, which needs to be keyed in, is the folder name. The folder name is the virtual folder name in the media server. After all the information is entered in, the streaming process can be initiated by clicking on the play button of the streaming media player. The streaming media player streamed the multimedia file from the host computer as shown in Figure 14.

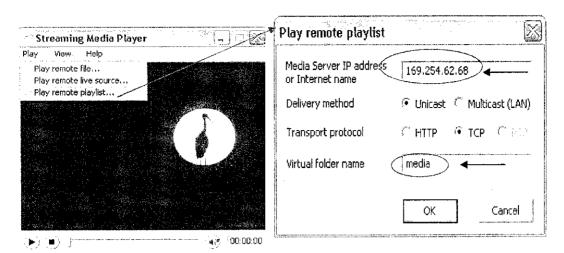


Figure 13: The streaming media player and the configuration.



Figure 14: The streaming media player streamed from the host computer.

## CHAPTER 4 RESULTS AND DISCUSSION

### 4.1 Experiment 1: Ad Hoc (Peer-to-peer) Wireless Network Streaming

 An ad hoc wireless network is set up between Laptop X (Intel Centrino<sup>®</sup>) and Desktop A (USB wireless adapter) in the laboratory at EE block 22 as shown in Figure 15. (*Please refer to section 3.4 for setting up wireless network*).

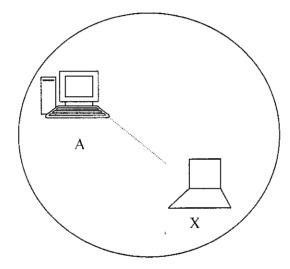


Figure 15: Ad-hoc Wireless Local Area Network

2. The streaming process as mentioned in section 3.5 is done within the wireless network created where Desktop A act as the media server while Laptop X streamed from the server.

The multimedia file streaming in ad hoc wireless network is successful as Laptop X (client) can stream the multimedia file smoothly from the Desktop A (host) provided Laptop X is within the coverage range of Desktop A. When the Laptop X is moved around within the coverage range, the two nodes stay connected and they are able to stream the multimedia files. The performance of the streaming process depends on the distance of node from the host. The distance between Desktop A and Laptop X is 1.80m and the data transfer rate is 5714 Kbit/s.

## 4.2 Experiment 2: Three Nodes Multi-hop Wireless Network Streaming

- An ad hoc wireless network is set up between Laptop X (Intel Centrino<sup>®</sup>) and Desktop A (USB wireless adapter) in the laboratory at EE block 22 as shown in Figure 15.
- 2. The Laptop X is moved far away from the Desktop A until the wireless connection no longer exists. (See Figure 16)

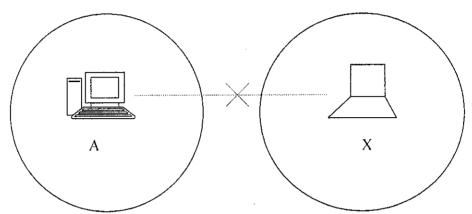


Figure 16: Two different networks with independent coverage range

 An additional node Laptop Y (Intel Centrino<sup>®</sup>) is added in the middle of Desktop A and Laptop X as shown in Figure 17.

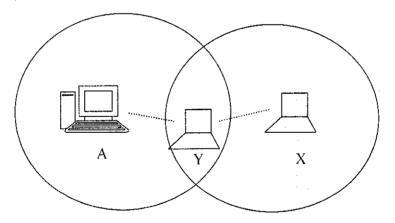


Figure 17: A multi-hop wireless local area network

4. The streaming process as mentioned in section 3.5 is done within the multihop wireless network created where Desktop A act as the media server while Laptop X and Laptop Y streamed from the server. When Laptop X is moved out of the coverage range of Desktop A, the wireless connection no longer exists. Therefore, the streaming process fails to perform and both nodes have independent network with two different coverage ranges as shown in Figure 16. By adding in an additional node – Laptop Y, in the middle of Desktop A and Laptop X as shown in Figure 17, the Laptop Y acts as a "bridge" for Desktop A and Laptop X. All nodes (A, X and Y) stay connected together under the same network and the streaming process is successful within this network where Desktop A acts as the server while Laptop X and Y as receiving nodes. The distance between Desktop A and Laptop Y is 1.80m and the data transfer rate is 5714 Kbit/s while the distance between the receiving end node Laptop X and Desktop A is 15.78m and the data transfer rate is 3517 Kbit/s. This shows that the distance does affect the performance of multimedia streaming in WLAN.

### 4.3 Experiment 3: Four Nodes Multi-hop Wireless Network Streaming

- 1. A multi-hop wireless network is set up between Desktop A (USB wireless adapter), Laptop X and Laptop Y (Intel Centrino<sup>®</sup>) as shown in Figure 17.
- 2. An additional node Desktop B (USB wireless adapter) is added at the far end of the network as shown in Figure 18.

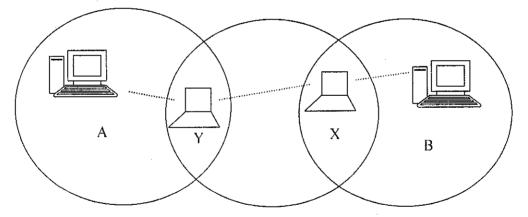
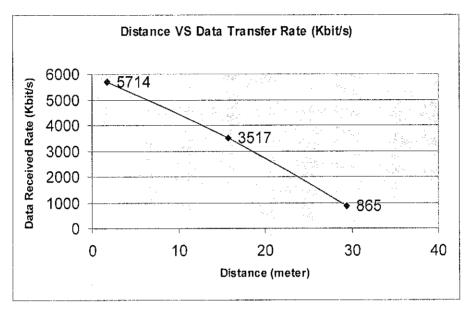


Figure 18: Multi-hop wireless local area network with four nodes

- 3. The streaming process as mentioned in section 3.5 is done within the multihop wireless network created where Desktop A act as the media server while Laptop X, Laptop Y, and Desktop B streamed from the server.
- 4. The distance between the host and the clients is measured.
- 5. The data transfer rate of every node is recorded and analyzed.

	A to B (far)	A to X (medium)	A to Y (near)
Distance (m)	29.38	15.78	1.80
Total packets sent	7988	7865	7924
Total time (s)	148	136	128
Data sent rate (packet/s)	53.97	57.83	61.91
Data received rate (Kbit/s)	865	3517	5714

Table 1: The measurements of the data transfer rate in multi-hop wireless network



*Figure 19: The distance VS data transfer rate (Kbit/s)* 

Desktop A and Desktop B are in two different networks, if Laptop X and Laptop Y are not placed in between. With Laptop X and Laptop Y acting as bridges, all nodes A, B, X, and Y are connected under the same network. It is found that the multimedia file can be successfully streamed between Desktop A and Desktop B.

The data transfer rate for different distances are investigated to evaluate the performance of the multi-hop wireless network. A two minutes video clip is streamed and the data transfer rate are observed. This is summarized in table 1. Figure 19 shows the graph is plotted from the result in table 1 and it shows that the distance of the node from the host does affect the data received rate at the clients. The nearer the node toward the host, the higher the data transfer rate and vise versa.

## 4.4 Experiment 4: Multi-hop Wireless Network Streaming without Access Point

- 1. A wireless network is created by Laptop Y (Intel Centrino<sup>®</sup>).
- 2. Laptop Y is moved near to Desktop A (USB wireless adapter) and Desktop A establishes the connection to the wireless network.
- After Desktop A is connected, Laptop Y is moved to the location of Desktop
   B (USB wireless adapter), which is far away from Desktop A.
- 4. This will cause Desktop A to lose connection to the network because it is outside the coverage range of the wireless network. However, now, Desktop B is connected to the network. The scenario is clearly illustrated in Figure 20.

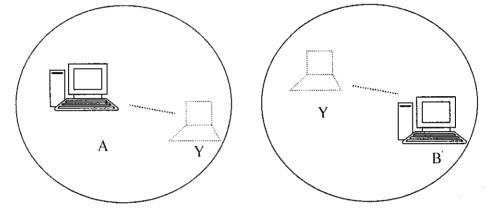


Figure 20: The connection of laptop Y to desktop A and desktop B

- 5. The Laptop Y is disconnected from the wireless network.
- 6. An additional node Laptop X (Intel Centrino<sup>®</sup>) is added in between Desktop A and Desktop B and a distance 1m away from the Desktop B as shown in Figure 21.

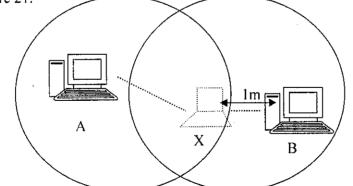


Figure 21: Multi-hop wireless local area network without access point.

- 7. The streaming process as mentioned in section 3.5 is done within the multihop wireless network created where Desktop A act as the media server while Laptop X and Desktop B streamed from the server.
- 8. The data transfer rate of the receiving node (Desktop B) is recorded in table 2.
- 9. The distance between the bridge (Laptop X) and the receiving node (Desktop B) is changed to 2m, 3m, 4, and 5m.
- 10. The data transfer rates of the receiving node for every distance are recorded in table 2.

 Table 2: The data transfer rate in multi-hop wireless network (moving bridge)

Distance between X and B	1m	2m	3m	4m	5m
Total packet received	· · · · · -	14568	13856	14540	15238
Total time elapsed (s)	-	367	198	426	498
Received rate (packets/s)	-	39.69	69.98	34.13	30.60
Received rate (Kbits)		531	1157	438	379

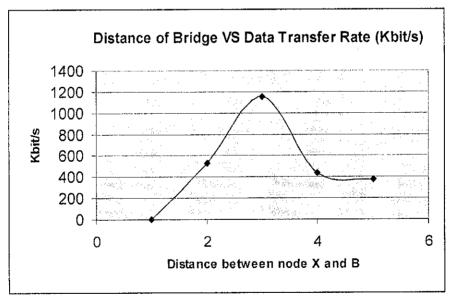


Figure 22: The distance of bridge from the receiving end node VS data transfer rate (Kbit/s)

The set up of the wireless network by using another node - Laptop Y, and disconnect it later on, is to provide Desktop A and Desktop B two independent wireless network. An additional node – Laptop X is then placed in between Desktop A and Desktop B as a bridge. The usage of the extra node Laptop Y to create the network is to avoid the existence of the access point where the Laptop Y is the originator of the wireless network, and thus acts as an access point. If this is the case then the wireless network created is not a purely ad hoc network, as it still needs a central main node for the network to exist.

The Laptop X is placed in between Desktop A and Desktop B as a bridge to connect all three nodes under the same wireless network. The streaming process is successfully done from Desktop A to Desktop B. This shows that the wireless network created is purely ad hoc and no access point is needed.

The performance of the streaming is observed when the bridge is placed at different distance from the receiving node and the result is summarized in table 2. It is found that streaming process is failed to perform when the bridge is placed about 1 meter from the receiving node. This is probably caused by the bridge is too far away from the host as it totally lost connection from the wireless network. The performance of the streaming process is optimum when the distance between the bridge and the receiving node is at 3 meters. The performance of streaming process degrades when the bridge is moved further than 3 meter away from the receiving node. The scenario is clearly illustrated in Figure 22.

## CHAPTER 5 CONCLUSION

The advantage of the streaming process allows the users to download and playback the multimedia file simultaneously. Thus, no big capacity of disk space required for storing an unknown content multimedia file. Streaming can be used for multimedia distribution without users having the illegal copies of the file. It allows the owner of the file to share the contents but then user would not be able to own the file.

The multimedia file must be streamed via two or more Wi-Fi enabled computer without passing through any access point. A node that is totally outside the coverage range of the network can also stream the multimedia files by using multi-hop wireless network passing through a bridge. The distance of the nodes from the host does affect the data transfer rate and thus affecting the performance of the streaming. The nearer the client node to the host, the higher the data transfer rate and vice versa. Every device has different signal strength and coverage range. The experiment results can be more accurate if all the wireless devices used are of the same model. The distance between the bridge and the receiving end node is about 3 meters to have the best performance of streaming process in multi-hop wireless network.

The streaming process is improved with no limitation of nodes. Therefore, a node outside the range can also stream by multi-hop networking. The signal strength of the wireless network can be improved to do larger distance multimedia files streaming. To improve this project, it is proposed that software will allow the users to choose any of the multimedia files in the playlist to be streamed rather than just playing all the multimedia files in order.

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## APPENDICES

## The results of packets transfer in USB wireless adapter

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66 38.341524 67 38.341690	169.254.45.10 169.254.45.10	225.0.0.1	UDP	Source por	t; 1254	Destination	n port: 5001
68 38,463496 69 38,553030	169.254.45.10 169.254.45.10	169.254.141.176 225.0.0.1	TCP UDP :-	<ul> <li>5119 &gt; 104</li> <li>Source pot</li> </ul>			7 win=17504 Len=0 n port: 5001
70 38, 553258	169.254.45.10	225,0.0.1	UDP	Source po	rt: 1254	Destination	n port: 5001
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73 38.599321	169.254.45.10	225.0.0.1	UDP : UDP :	Source pol	τ: 1254	. Destination	n port: 5001 n port: 5001
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79:38.818116 80:38.818256	169.254.45.10 169.254.45.10	225.0.0.1 225.0.0.1	UDP .	Source po	nt:∬1254	Destination	n port: 5001 n port: 5001
81 38.904762	169,254.45.10	169.254.141.176	TCP	5119 > 10	15 [PSH,	ACK] Seg-1 A	Ack=17 win=17504 Len=311
82 38,985920 83 38,986102	169.254.45.10 169.254.45.10	225.0.0.1 225.0.0.1	UDP UDP	Source po	rt: 1254	Destination Destination	n port: 5001 =17 Win=17504 Len=1460
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89 39.031204	169,254.45.10	225.0.0.1	UDP	Source po	rt: 1254	Destinatio	n port: 5001
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42 40 261550 43 40 261597	169.254.45.10	169.254.141.176	TCP	5119 > 10	45 [ACK]	Seq=19004 A	ck=17 win=17504 Len=1460 ck=17 win=17504 Len=1460
44 40.261670 45-40.262542	169.254.45.10 169.254.141.176	169.254.141.176 169.254.45.10	TCP TCP	TCP wind	ow Updat	e 1045 > 51	19 [ACK] Seq=17 Ack=1900
46 40,262555	169.254.45.10	169.254.141.176 169.254.141.176	TCP	5119 > 10 5119 > 10	45 [PSH, 15 [ACK]	ACK] Seg=21 Seg=23384 4	924 Ack=17 win=17504 Len ck=17 win=17504 Len=1460
147 40.262575 148 40.262591	169.254.45.10 169.254.45.10	169,254,141,176	TCP	5119 > 10	45 (РЗН,	ACK] Seq=24	844 ACK=17 WIN=17504 LBN
149 40.262607 150 40.262622	169.254.45.10 169.254.45.10	169.254.141.176 169.254.141.176	TCP	5119 > 10 5119 > 10	45 [ACK] 45 [ACK]	Seq=20304 A Seq=27764 A	ck=17 Win≃17504 Len=1460 ck=17 Win=17504 Len=1460
151 40.262642	169,254.45.10	169.254.141.176	TCP	5119 > 10	45 [ACK]	Sed=29224 A	ck=17 win≕17504 Len=1460 ck=17 win≕17504 Len=1460
LS2 40.262667 LS3 40.262685	169.254.45.10 169.254.45.10	169,254.141.176 169,254,141,176	TCP TCP	5119 > 10	45 [ACK]	Seq=32144 A	ck=17 win=17504 Len=1460
154 40.264647	169.254.45 <b>.1</b> 0	225.0.0.1	UDP TCP			Destinatio seg=17 Ack+	n port: 5001 21924 Win=17520 Len=0
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159 40.269513	169.254.141.176	169,254.45.10	TCP .	1045 > 51	19 [ACK]	Seq≕17 Ack=	23384 Win=17520 Len=0 ck≈17 Win=17504 Len=1460
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16 179, 712225	169.254.141.176	169.254.45.10	TCP	1048 > 5119 [ACK]	Seg=17 Ack=1126104 Win=17520 Len=0
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	169.254.141.176	169.254.45.10	TCP		Seg=17 Ack=2456015 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP		Seq=2470615 Ack=17 Win=17504 Len=146
	169,254,45,10	169.254 141.176	TCP		5119 > 1048 [ACK] Seg=2472075 Ack=17
	169.254.45.10	225.0.0.1	UDP	Source port: 1254	
	169.254.45.10	225.0.0.1	UDP	Source port: 1254	Destination port; 5001
54 182,073029	169.254.141.176	169.254.45.10	TCP	1048 > 5119 [ACK]	Seq=17 Ack=2457475 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP		5119 > 1048 [ACK] Seq=2473535 Ack=17
	169.254.141.176	169.254.45.10	тср		Seg=17 Ack=2460395 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP		ACK] 5eq=2474995 Ack=17 win=17504 Le
	169.254.45.10	169.254.141.176	TCP		5119 > 1048 [ACK] Seq=2476455 Ack=17
	169.254.141.176	169.254.45.10	TCP		Seq=17 Ack=2463315 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP		Seq=2477915 Ack≠17 Win=17504 Len=146 5119 > 1048 [ACK] Seq=2479375 Ack≠17
	169.254.45.10 169.254.141.176	169.254.141.176 169.254.45.10	TCP		Seg=17 Ack=2466235 Win+17520 Len=0
	169.254.45.10	169.254.141.176	TCP		Seg=2480835 Ack=17 Win=17504 Len=146
55 162.079971 Sa 292 0200004	169.234.43.10	169.254.141.176	TEP -		5119 > 1048 [ACK] SEG=2482295 ACK=17
65 182 085835	169.254.45.10 169.254.141.176	169.254.45.10	TCP		Seg=17 Ack=2469155 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP		seg=2483755 Ack=17 win=17504 Len=146
		169.254.141.176	TCP	TCP Window Full	5119 > 1048 [ACK] Seg=2485215 Ack=17
	169.254.141.176	169.254.45.10	TCP	1048 > 5119 [ACK]	Seg=17 Ack=2472075 Win=17520 Len=0
	169.254.45.10	169.254.141.176	TCP	5119 > 1048 [ACK]	Seg=2486675 Ack=17 Win=17504 Len=146
	169.254.45.10	169.214.141.176	TCP		5119 > 1048 [ACK] Seq=2488135 Ack=17

The results of packets transfer in Ir	ntel Centrino <sup>®</sup> laptop
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	। প্রান্থ প্রান্থ প্রান্থ প্রান্থ		<i>تيبية</i>	<u>(11</u> )	~	•4	لنتيب			Ехрге			÷		هـ ب	~~	~~~	~		- 	<u>ب</u> ،		0.	
Edter:						· · · ·	the other						- The	• <b>7</b>										
No 5812	Time 99 <b>.1</b> 75973	5ourc 1.69	. 254 .	45.10	2	22	stination	2.1		u	otocol DP	Sour	ce j	port:	1254	Des	tinat	ion	port	: 500	1			
5813 5814	99.176817 99.256537	169 . 169	.254. ,254.	45.10 45.10	2 . 2 .	22	5.0.1	).1 ).1		00	DP DP ·	- Sour - Sour	CE	port:	1254	Des Des Des	tina	100	port	: 500	1			
.581.5	99.256537 99.256733 99.257395	169	, 254.	141.3	L76	1.6	9.25	4.45.	10	$\sim T_{c}$	- P	: 1050		5119.	IFIN.		260	•U AC	K≓ZBI	UJ 124	ਂਆ ⊓ਸ=	16528	3 Lên	=0
5810	99,259089	T 0 9	. 2 . 4 .	-4 J - T (	,	16	9.25	1.141	.176	TC	2P	5119		1050	[ACK]	Des Seq= Des	2803:	24 A	ck=1	win=	17504	Len	=0	
5818	99,303889	169	.254.	45.10	)	22	5.0.1	0.1		U) Trent	DP mpsalst	Sour	ce	port:	1254 TETN	Des ACK]	tinat	:10n -2803	port 124 -	: 500 Arka1	1 ∵win≞	17502	Piten	20
2850	99.304272 99.304291	169	.254.	45.1	L76			1.45.		ΤC	CP	1050	>	5119	[ACK]	_ Seq≠	1 Aci	(=280	3125	win-	16528	Len	=0	- 5
5821	99.345783	169	.254.	45.10	0		5.0.1	0.1 1.45.:	10	UC TC	0P	Sour	ce j	port: 5119.	: 1254 [SVN]	Des Seğ=	tina: Oʻaci	:10n (=0. W	port in≕u	: 500 6384	1 LenãC	MSS	=1460	
5823	99,368270 99,370354	169	.254.	45.10	) · · ·	16	9,25	1.141	176	े ें ज	CP ()?	5119	158-1	1055	ÎSYN.	ACK]	. Seq.	=0 Ac	k=1 1	whin=1	7520	Len=(	D, MSS	-1
	99.370388 99.370454			141.1				1.45. 1.45.:		т т Т		1055	>	511.9	FPSH.	Seq= ACK]	Sea	-1 AC	k=1 %	w1n=1	7520	Len=3	16	
5826	99.492415	169	. 254.	45.1	0	16	9.254	1.141	.176	ΤC	ΞP	5119	>	1055	[ACK]	Seq= A⊂K] Des		-17	Win=:	17504	Lene	0		
	99,589918 99,591671			45.1				4.141 0.1				SOUP	ce j	port:	LPSB 1.254	Des	seq: t1na	∈⊥ AC tion	port	: 500	1/504	Len	- 2 2 2	
5829	99.592599	169	.254.	45.10	p -	22	5.0.1	0.1	1.1	- U(	DP ·	Sour	ce i	port:	1254	. Des	tina	cion –	port	: 500	1			с. с
	99,686422 99,686488	169	.254.	45.1	176	16	9.25	4.141 4.45.	.176 10	Т( Т(		1055	> .   >	5119	[ACK]	ACK] Seq=	5eq 17 A	=324 ck≈17	84 W	in=17	'520 L	.en=0	11-14	00
5832	99,687968	169	.254.	45.1	p i	· 22	5.0.0	0.1		- UI		Sour	ce	port:	: 1254	Des	tina	tion	port	: 500	1			
5834	99,688920 99.690250	169	.254.	45.1	0	22	5.0.1	0.1		UC UC		Sour	cei	DOPT :	1254	Des	tina	tion	dort	: 500	1			
5835	99.713697	169	. 2'54 .	45.1	D	22	5.0.1	0.1	170	UC UC TC	DP	Sour	ce	port	(1254 [ACK]	Des Seg=	tina: 1794	tion Ack-	port	: 500 iow17	1.	an=1/	160	
	99.741489 99.742935	169	.254.	45.1	0			4.141				5119		1055	[ACK]	Seq= Seq=	3244	Acka	17 W	in-17	504 L	en=14	160	
5838	99.742968	169	.254,	141.1	176	16	9.25	4.45.3	10	T		5110		1055	- FACKI	500-	4704	ACk a	17 w	in=17	504 1	Pn≃14	160	
	99.744364 99.747368	169	.254.	45.1	Ď	16	9.25	4.141 4.141	176	ΤC	CP	5119	> > :	1055	[ACK]	Sed=	6164	ACk =	17 W	in⊨17	'504 L	en=14	160	
	99.747383	169	,254.	141.1 45.1	176			4.45.1 4.141			CP CP	1055	i > .	5119	[A⊂K]	Seq=	17 A	CK=/0	24 W	1/1=1/	520 L	en≖u		
5843	99.749088 99.751043	169	.254.	45.1	0	16	9.25	4.141	.176	т	CP	\$119	) > C	1055	<b>JACK</b>	Sed=	9084	Ack=	17 w	in=17	504 L	.en=14	160	
5844	99.751056 99.752498	169	.254.	141.3	176 1			4.45. 4.141			CP CP	1055	i > i > :	5119 1055	LACK	Seq=	17 A 1054-	ck =10 4 Ack	=17	wnn≕1 Win=1	7504	Len=1	J 1460	
5846	99.753958	169	.254.	45.1	0	16	9,25	4.141	.176	Τ¢	CP	5119	)>:	1055	[ACK]	Seq=	1200	4 Ack	-17 '	Win=1	7504	Len=1	1460	
	99,753973 99,755666			141.: 45.1				4.45.° 4.141		TC TC		1055	) > ) > )	5119 1055	[PSH	Seq= ACK]	Sea	ск=13 =1346	404 4 AC	win≞⊥ k≑17	./s∠u win=1	.7504	Len≓	14
5849	99.755693	169	.254.	141.:	176	1,6	9.25	4.45.3	10	Τ¢	CP	1055	i: >	5119	[ACK]	Sea+	17 A	ck =14	924	win=1	7520	Len=(	0	
	99.374793 99.874810			.45 <b>.1</b> 141.:				4.141 4.45		T (		5119	> .   >	1055 5119	EACK	ACK] Seg=	5eq 17 A	=1981 ck=21	5 AC	κ∝ı/ Win=1	wnn=1 7520	.7504 Len≐(	Len≃ 0	14
5859	99.875204	169	.254.	45.1	Ô.	16	9.25	4.141	.176	τ¢	CP	5119	) > :	1055	[PSH,	ACK]	seq	=2127	'G AC	k≂17	win=1	7504	Lene	40
	99.920772 99.921733			45.1			5.0.	0.1 0.1		UI · UI		Sour	ice i	port: port:	1254	. Des	tina tina	tion	port	: 500	1			
5862	99.923130	169	.254.	45.1	0	22	5.0.	0.1		U	DP	Sour	ce	port:	: 1254	Des Des	tina	tion	port	: 500	1			e i
	99,923981 99,925514	169	.254.	45.1	0. 0	16	5.0/	U.1 4,141	.176	UI	CP	5119	1 > 1	1055	[PSH	ACK ]	seq	-2131	.6 AC	k=17	win=1	7504	Len=	
5865	99.925541	169	. 254.	.141.1	176	16	9.25	4.45.	10	T	CP	1055	12.	5119	[ACK]	Seq=	17 A	ck=22 -2272	776	W1n=1 2−17	.7520 win-1	Len=0 7504	) ien	10
	99.926125 99.927906	169	.254.	45.1	ů Č	16		4,141 4.141				5119	5	1055	[PSH,	ACK] ACK]	Seq	-2297	S AC	k=17	Win=1	7504	Len=	14
5868	99.927920	169	.254.	141.	176	16	9.25	4,45. 4,141	$\frac{10}{176}$	T T		1055	) <b>&gt;</b>	5119	LACK.	Seq= ACK]	1/A	ÇK =24	455	Mu⊔≞T	./520	Len=	0	
5870	99.928263 99.965365	169	.254.	45.1 45.1	ŏ	22	5.0	0.1	. 1 . 0	U	DP	Sour	°¢e	port	: 1254	Des	tina	tion	port	: 500	1			Ξ.
	99.966127 100.00566	169	. 254	45.1	0 ·	22	5.0.	0.1 4.141	176		DP	Sour \$110	-CE ) > '	port: 1055	: 1254 (РSH.	Des ACK]	tina Seo	tion =2445	2 AC	: 500 k=17	n Win≕1	.7504	Len=	14
5873	100.00568	169 (	.254.	.141.:	176	16	9.25	4.45.	10	T	CP	1055	i >	5119	[ACK	Seq-	17 A	ck ≠25	912	win=1	7520	Lenet	0	
	100.00666	L 169 7 169	254	45.1	0		9.25	4,141 0.1	,176	. те		Saur	) . 'ce	1055 Dort:	(РSН, : 1254	ACK] Des	tina	=2091 tion	2 AC port	κ≃⊥/ : 500	wnn≂⊥ )1	.7304	Lens	
5876	100.00950	5 169	. 254 .	45.1	0	22	5.0.	0.1		U)	OP .	Sour	ce	port	: 1254	Des	tina	tion	port	: 500	)1	7504		
	100.01068	$2 169 \\ 6 169$	.254	45.1	0 176	16	9.25	4.141 4.45.	$\frac{176}{10}$	Τ· Τ·		1055	i >	5119	1 ACK	ACK] Seq	17 A	CK = 2 /	275	win=1	.61.57	Len=	0	
5879	100.04926	0.169	.254.	45.1	0	22	5.0.	0.1	÷ 1		DP .	Sour	ce	port	1254	Des Des	tina	tion	port	: 500	1		1.1	÷.
	100.05009			45.1		16	5.0. 9.25	4.141	.176	-: U ⊺(	CP	5119	) > .	1055	[PSH.	ACK]	seq	=2727	'5 AC	k=17	W1n=1	7504	Len=	84
5882	100.09434	7 169	.254.	45.1	ο.	22	5.0.	9.1	5 A.	·∘ U	DP					Des								1
5972	101,40575	7 169 0 169	, 254. . 254	45.1	0.	22	5.0.	0,1 0.1		t de la	DP DP	Sour	'ce 'ce	port: port:	: 1754	Des Des	tina	tion.	port	: 500	1.		1.1	Ĵ
5974	101.43136	5 169	.254.	45.1	0	16	9.25	4.141	.176	<b>T</b> 1	CP	5119	3 >	1055	[PSH	ACK]	seq	=6146	j1 A⊂	k=17	win=3			14
	101.43138							4.45. 4.141			CP CP	5119	3 > 1	1016	LACK.	_Seq=   Seq=	5787	382 A	ck=1	7 Wir	1 <b>=175</b>	)4 Lei	n <b>≈14</b> 6	0
13519	124.44513	7 169	.254.	.45.1	0	16	i9.25	4,141	.176	Ť	CP	5119	3 >	1056	[ACK]	Seq-	5788	842 A	.ck=1	.7 w1r	1=1750	)4 Lei	n=146	0
13520	124.44515 124.44894	2 169 1 169		.141.: .45.1				4.45. 4.141			CP CP	5119	) > } >	1056	LACK	Seq=	5790	302 A	<pre>s030 <ck=1< pre=""></ck=1<></pre>	2 w1r 7 w1r	1=1750	0 Lei )4 Lei	n=146	jQ
13522	124.45067	5 169	.254.	45.1	0	16	9.25	4.141	.176	T	CP	5119	) > I	1056	[PSH,	, ACK]	Seq	≃5791	.762	Ack=1	.7 W I	1=175	04 Le	n-
13523	124.45070 124.45241			.141. .45.1				4.45. 4.141			CP CP	5119	) > } >	1056	ACK	Seq=	5793	222 A	9522 ck=1	zwir 7wir	1=1750	0 Lei )4 Lei	n=146	0
13525	124.45441	6 169	1.254.	.45.1	Q	10	9.25	4.141	.176	т	CP CP	5110	1.5	1056	[ACK]	Seq=	5794	682 /	∖ck=1	7 Wir	1=175(	)4' Lei	n=146	0
	124.45442 124.45592			.141. .45.1		10	9.25	4.45. 4.141	176		CERPC	. Alte	ric	onte	xt.c	all ic	1: 67	11116	8.4	<ul> <li>cont</li> </ul>	ext '	items	. lst	
13528	124.45765	9 169	254	,45.1 .141.	0	19	19.25	4.141 4.45.	.176	T	CP CP	5119	9 > 1 \	1056 5119	[ACK]	Seq=	-5797 17 A	602 4 ck=57	.ck=1	.7 Wir 2 Wir	1=175( 1=1752	)4 Lei 20 Lei	n⊨14€ n=0	0
	124.45767 124.45932	0 169	.254	,45.1	0	10	i9.25	4,141	,176	Ť	CP	5119	•	1056	[ACK]	) Seq-	5799	Q62 A	∖ck=1	7 Wir	1=175(	)4 Lei	n=146	0
13531	124.46086	9 169	1.254	.45.1 .141.	0	10	59.25	4.141 4.45.	.176	T	CP CP	5119 1056	3 > 5 >	1056 511 9	[ACK]	Seq Seq	5800 17 A	>22 4 ck ⇒ 58	ACK=1 301.98	./ Wir 2 Wir	า≕175( 1≕1752	24 Lei 20 Lei	n≕146 n≕0	Û
13533	124.46288	2 169	1,254	.45.1	0	16	9.25	4.141	.176	T	CP	5119	а х.	1056	EACK	l Sena	5801	982 A	uck=1	7 Wii	1=1750	)4 Le	n= <b>14</b> 6	Û
13534	124.46442 124.46446	4 169	254	.45.1	0	16	9,25 9,25	4,141 4,45,	.176 10	Ϋ́ Ti	CP CP	5119 1056	9 > 5 >	1056 5119	LACK	Seq Seq	-5803 17 A	442 A ⊆k=58	ACK=1 30490	.7 Wir 2 Wir	n=175( n=1752	20 Le	n=146 n=0	Û
13536	124.46603	5 169	,254	.45.1	0	22	25.0.	0.1	T.A.	U	DP .	sour	°ce –	port	125	1 Des	πına	ເາດຄ	port	: 500	л.			
13537	124.46724	0 169	254	.45.1	Ú		5.0. 5.0.				IDP IDP	Sour	-ce	port port	125 125	Des Des	tina tina	tion.	port	: 500 : 500	л Л			1
13539	124.47005	5 169	254	.45.1	0	16	59.25	1.141		۳ i	CP	5119	€ >	1056	[А⊂К]	] seq∍	5804	902 A	∖ck=1	.7 ∀1ı	1= <b>17</b> 5(	04 Le	n=146	0
13540 13541	124.47175	2 169 4 169		.45.1 . <b>1</b> 41.				4.141 4.45.			CP CP				[ACK]	] Seq=	17 A	ck⇒58	30782	2 1/1	1=175( 1=175)	20 Le	n=0	
13542	124.47322	5 169	.254	.45.1	0	10	59.25	4.141	.,176	т т	CP	5119	€ €	1056	[ACK [ACK	] Seq.	5807	822 /	∖ck=1	.7 Wi∣	1-1750	04 Le	n≂146 n=146	0
	124.47524							4.141			CP													

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	74393 446.168946	169.254.45.10	169.254.141.176	TDS	Unknown Packet Type: 0 (Not last buffer) (Not last buffer) (1
	74394 446.168958	169.254.141.176	169.254.45.10	TCP	1064 > 5119 [ACK] Seq=18 Ack+17385431 Win=17520 Len=0
	74395 446.170467	169.254.45.10	169,254,141,176	TCP	[TCP segment of a reassembled PDU]
	74396 446.172134	169.254.45.10	169.254.141.176	TCP	[TCP segment of a reassembled PDU]
	74397 446.172148	169.254.141.176	169.254.45.10	TCP	1064 > 5119 [ACK] Seq=18 Ack=17389351 Win=17520 Len=0
	74398 446.175612	169,254,45,10	169.254,141.176	TCP	[TCP segment of a reassembled PDU]
	74399 446.177277	169,254,45.10	169.254.141.176	TCP	[T⊂P segment of a reassembled PDU]
	74400 446.177296	169.254.141.176	169,254,45.10	TCP	1064 > 5119 [ACK] Seq=18 Ack=17391271 win=17520 Len=0
	74401 446,178690	169,254,45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74402 446,179865	169.254.45.10	169.254.141.176	TCP	[TCP segment of a reassembled POU]
	74403 446.180772	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74404 446,185147	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
1	74405 446.187934	169.254.45.10	169.254.141.176	TCP	[TCP segment of a reassembled PDU]
	74406 446.187955	169,254.141.176	169.254.45.10	TCP	1064 > 5119 [АСК] Seq=18 Ack=17392263 Win=16528 Len=0
	74407 446,227637	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74408 446,228531	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74409 446,419264	169,254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74410 446.420004	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74411 446.437135	169.254.45.10	225,0.0.1	UDP	Source port: 1254 Destination port: 5001
	74412 446,438037	169.254.45.10	225.0.0.1	UDP	source port: 1254 Destination port: 5001
	74413 446,506178	0.0.0.0	255,255,255,255	DHCP	DHCP Discover - Transaction ID 0x1a471868
	74414 446 539108	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74415 446.539844	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74416 446.550491	169.254.45.10	225,0.0.1	UDP	Source port: 1254 Destination port: 5001
	74417 446.551374	169.254.45.10	225.0.0.1	UDP	Source port: 1254 Destination port: 5001
	74418 446,746956	169.254.45.10	225,0.0.1	UDP	Source port: 1254 Destination port: S001
4	74/10 4/6 757797	160 754 45 10	775 0 0 1	IIND	Source hort 1754 Destination port. 5001
		,	Timeren OA. Th	AC.	ED lanton nousilth

Figure 24: The ACER laptop results

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